# Goncharov's programme, and depth reductions of multiple polylogarithms

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# Multiple polylogarithm definitions

#### Definition

Multiple polylogarithm (MPL) is

$$\operatorname{Li}_{k_1, k_2, \dots, k_d}(x_1, x_2, \dots, x_d) \coloneqq \sum_{0 < m_1 < m_2 < \dots < m_d} \frac{x_1^{m_1} \cdots x_d^{m_d}}{m_1^{k_1} \cdots m_d^{k_d}}, \quad |x_i| < 1.$$

Generalises

$$\text{Li}_1(x) = -\log(1-x) = \sum_{m=1}^{\infty} \frac{x^m}{m}.$$

- Weight is  $n = k_1 + \cdots + k_d$ ,
- $\blacksquare$  Depth is d.

Goal: Fundamental challenge around MPL's is to understand functional equations, identities, and the filtration by depth.

# Functional equations for polylogarithms

■ Functional equations for Li<sub>1</sub>

$$\log(x^{-1}) = -\log(x)$$
,  $\log(xy) = \log(x) + \log(y)$ .

■ Propagation to dilogarithm via  $\frac{d}{dx} \operatorname{Li}_2(x) = \frac{1}{x} \operatorname{Li}_1(x)$ ?

#### Proposition (Dilogarithm symmetries)

Dilogarithm satisfies two symmetries 
$$\begin{aligned} &\underset{=\mathrm{Li}_2(1)=\frac{\pi^2}{6}}{\mathrm{Li}_2(x)+\mathrm{Li}_2(1-x)} &= \overbrace{\zeta(2)} - \log(1-x)\log(x)\,, \\ &\mathrm{Li}_2(x)+\mathrm{Li}_2(x^{-1}) &= -\zeta(2) - \frac{1}{2}\log^2(-x)\,, \quad x \not\in [0,1] \end{aligned}$$

#### Theorem (5-term relation, Abel, Spence, Kummer, ...)

For |x| + |y| < 1 we have 'fundamental identity'

$$\operatorname{Li}_{2}(x) + \operatorname{Li}_{2}(y) - \operatorname{Li}_{2}\left(\frac{x}{1-y}\right) - \operatorname{Li}_{2}\left(\frac{y}{1-x}\right) + \operatorname{Li}_{2}\left(\frac{xy}{(1-x)(1-y)}\right)$$
$$= -\log(1-x)\log(1-y)$$

# Relating different depths

Functions of different depths can be related:

$$\operatorname{Li}_{1}(x)\operatorname{Li}_{1}(y) = \operatorname{Li}_{1,1}(x,y) + \operatorname{Li}_{1,1}(y,x) + \operatorname{Li}_{2}(xy)$$

$$\leftrightarrow \quad \sum_{m=1}^{\infty} \sum_{\ell=1}^{\infty} \frac{x^m y^{\ell}}{m \cdot \ell} = \left( \sum_{m < \ell} + \sum_{m > \ell} + \sum_{m=\ell} \right) \frac{x^m y^{\ell}}{m \cdot \ell}$$

So-called stuffle product of MPL's  $\rightsquigarrow$  this generalises easily

More surprising when the depth reduces

#### Proposition (Lewin<sup>(?)</sup>, popularised by Goncharov, Zagier)

As a power-series identity, for |xy| < 1, |y| < 1,

$$\operatorname{Li}_{1,1}(x,y) = \operatorname{Li}_2\left(\frac{y(x-1)}{1-y}\right) - \operatorname{Li}_2\left(\frac{-y}{1-y}\right) - \operatorname{Li}_2(xy).$$

# Further depth reductions?

#### Proposition (Li<sub>1,1,1</sub> is depth 1, Lewin<sup>(?)</sup>, conceptionally by Goncharov, Zhao, ...)

$$\begin{array}{rcl} {\rm Li}_{1,1,1}(x,y,z) & = & - {\rm Li}_3\Big(\frac{1-xyz}{1-x}\Big) - {\rm Li}_3\Big(\frac{1-xyz}{xy(1-z)}\Big) + {\rm Li}_3\Big(\frac{(y-1)(1-xyz)}{(1-x)y(1-z)}\Big) + {\rm Li}_3(xy) \\ & & - {\rm Li}_3\Big(\frac{y(1-x)}{y-1}\Big) + {\rm Li}_3(1-x) - {\rm Li}_3\Big(\frac{y(1-z)}{y-1}\Big) + {\rm Li}_3\Big(\frac{-y}{1-y}\Big) + \textit{products} \end{array}$$

Can weight 4 be reduced to depth 1? Apparently not (e.g. Wojtkowiak), but:

#### Proposition ( $Li_{3,1}$ relations, mod depth 1, Zagier, Gangl)

$$\begin{split} \operatorname{Li}_{3,1}\!\left(\frac{1-x}{y},y\right) + \operatorname{Li}_{3,1}\!\left(\frac{x}{y},y\right) &= -\frac{1}{2}\operatorname{Li}_4\!\left(\frac{(1-x)y}{x(1-y)}\right) - \frac{1}{2}\operatorname{Li}_4\!\left(\frac{xy}{(1-x)(1-y)}\right) + \frac{1}{2}\operatorname{Li}_4\!\left(\frac{(1-y)y}{(1-x)x}\right) \\ &- \operatorname{Li}_4\!\left(\frac{1-y}{1-x}\right) - \operatorname{Li}_4\!\left(\frac{1-y}{x}\right) + \operatorname{Li}_4\!\left(-\frac{y}{1-y}\right) \\ &- \operatorname{Li}_4\!\left(1-x\right) - \operatorname{Li}_4\!\left(x\right) + 2\operatorname{Li}_4\!\left(y\right) + \textit{products} \end{split}$$

Question: How to predict or understand, such reductions and obstructions?

#### Motivations & connections elsewhere

#### Values of zeta functions:

- Dedekind zeta function of a number field F is  $\zeta_F(s) = \sum_{I \neq 0 \subset \mathcal{O}_F} N(I)^{-s}$ .
- Analytic class number formula:  $\operatorname{Res}_{s=1} \zeta_F(s)$  via regulator  $R_F$  (& arithmetic data).
- **Z**agier: conjecture for  $\zeta_F(n)$  via  $\text{Li}_n$  (& 'higher units' in K-theory).

Example: 
$$\zeta_{\mathbb{Q}(\sqrt{-7})}(2) = \frac{4\pi^2}{21\sqrt{7}} \left( 2 \underbrace{D_2\left(\frac{1+\sqrt{-7}}{2}\right)}_{\substack{\text{single-valued version of Li}_2}} + D_2\left(\frac{-1+\sqrt{-7}}{4}\right) \right)$$

■ Goncharov relates  $\zeta_F(n)$  to high-depth "Grassmannian polylogarithm".

Task: How to reduce to depth 1?

Also: hyperbolic geometry, cluster algebras, high-energy physics, . . .

# Motivic multiple polylogarithms

Recall: Iterated integral

$$I_{(\gamma)}(a;x_1,\ldots,x_n;b) \coloneqq \int_{a < t_1 < \cdots < t_n < b} \frac{\mathrm{d}t_1}{t_1 - x_1} \wedge \frac{\mathrm{d}t_2}{t_2 - x_2} \wedge \cdots \wedge \frac{\mathrm{d}t_n}{t_n - x_n}.$$

By term-wise integration of geometric series (with  $X_j := \prod_{i=1}^j x_i$ ),

$$\operatorname{Li}_{k_1,k_2,\dots,k_d}(x_1,\dots,x_d) = (-1)^d I(0;1,\underbrace{0,\dots,0,X_1}_{k_1},\underbrace{0,\dots,0,X_2}_{k_2},\dots,X_{d-1},\underbrace{0,\dots,0;X_d}_{k_d}).$$

■ Goncharov: upgrade of I to framed mixed Tate motives  $I^{\mathfrak{u}}(a; x_1, \ldots, x_n; b)$ ,  $I^{\mathfrak{u}}$  lives in connected, weight graded Hopf algebra  $(\mathcal{A}_{\bullet}(\overline{\mathbb{Q}}), \Delta)$ 

Idea: Can keep track of differential form and integration chain separately

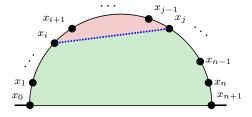
lacksquare Define  $\mathrm{Li}^{\mathfrak{u}}_{k_1,k_2,\ldots,k_d}(x_1,\ldots,x_d)$  via  $I^{\mathfrak{u}}$ 

### Lie coalgebra, cobracket, depth filtration

Write:  $\mathcal{L}_{\bullet} = \mathcal{A}_{\bullet}/\mathcal{A}_{>0}^2 = \text{motivic MPL's modulo products}$ 

■ Image of  $I^{\mathfrak{u}}, \operatorname{Li}^{\mathfrak{u}}_{k_{1},...,k_{J}}$  in  $\mathcal{L}_{\bullet}$  is  $I^{\mathcal{L}}, \operatorname{Li}^{\mathcal{L}}_{k_{1},...,k_{J}}$  ■ Obtain  $\Delta \colon \mathcal{L}_{\bullet}(F) \to \bigwedge^{2} \mathcal{L}_{\bullet}(F)$ 

$$\Delta I^{\mathcal{L}}(x_0; x_1, \dots, x_n; x_{n+1}) = \sum_{i < j} I^{\mathcal{L}}(x_i; x_{i+1}, \dots, x_{j-1}; x_j) \wedge I^{\mathcal{L}}(x_0; x_1, \dots, x_i, x_j, \dots, x_n; x_{n+1})$$



■ Depth filtration  $\mathcal{D}_d = \{ \text{depth} \leq d \text{ MPL's in } \mathcal{L}_{\bullet} \}$  has

$$\Delta(\mathcal{D}_d) \subset \mathcal{D}_d \wedge \mathcal{L}_1 \oplus \bigoplus_{\substack{i+j=d \ i,i \geq 1}} \mathcal{D}_i \wedge \mathcal{D}_j.$$

Why? Follows since  $I^{\mathcal{L}}(a; \overbrace{0, \dots, 0}; b) = \frac{1}{\ell!} I^{\mathcal{L}}(a; 0, b)^{\ell} \equiv 0 \pmod{\mathsf{products}}$ 

## Cobrackets and obstructions to depth 1

$$\Delta \operatorname{Li}_{4}^{\mathcal{L}}(x) = \log^{\mathcal{L}}(x) \wedge \operatorname{Li}_{3}^{\mathcal{L}}(x) \qquad \in \mathcal{L}_{1} \wedge \mathcal{L}_{3},$$

$$\Delta \operatorname{Li}_{3,1}^{\mathcal{L}}(x,y) = -\log^{\mathcal{L}}(x) \wedge \operatorname{Li}_{2,1}^{\mathcal{L}}(x,y) + \operatorname{Li}_{1}^{\mathcal{L}}(y) \wedge \operatorname{Li}_{3}^{\mathcal{L}}(x) - \operatorname{Li}_{1}^{\mathcal{L}}(y) \wedge \operatorname{Li}_{3}^{\mathcal{L}}(xy)$$

$$- \operatorname{Li}_{1}^{\mathcal{L}}(xy) \wedge \operatorname{Li}_{3}^{\mathcal{L}}(x) + \operatorname{Li}_{1}^{\mathcal{L}}(xy) \wedge \operatorname{Li}_{3}^{\mathcal{L}}(y)$$

$$+ \operatorname{Li}_{2}^{\mathcal{L}}(y) \wedge \operatorname{Li}_{2}^{\mathcal{L}}(xy)$$

$$\in \mathcal{L}_{1} \wedge \mathcal{L}_{3}$$

$$\in \mathcal{L}_{2} \wedge \mathcal{L}_{2}$$

lacksquare Write  $\overline{\Delta}$  to mean ignore  $\mathcal{L}_1$ , then we obtain

$$\overline{\Delta} \operatorname{Li}_{4}^{\mathcal{L}}(x) = 0, \qquad \overline{\Delta} \operatorname{Li}_{3,1}^{\mathcal{L}}(x,y) = \operatorname{Li}_{2}^{\mathcal{L}}(y) \wedge \operatorname{Li}_{2}^{\mathcal{L}}(xy).$$

Hence  $\operatorname{Li}_{3,1}^{\mathcal{L}}(x,y) \neq \sum \operatorname{Li}_{4}^{\mathcal{L}}$ 's.

Question: Is this the only obstruction?

#### Goncharov Depth Conjecture (depth 1, simplified version)

A linear combination of MPL's has depth 1 if and only if  $\overline{\Delta}$  vanishes.

### Expectations and evidence

Weight 2, 3: No space for non-trivial contribution, so

$$\overline{\Delta}(\mathcal{L}_2) = 0, \quad \overline{\Delta}(\mathcal{L}_3) = 0.$$

Expect  $\operatorname{Li}_{1,1}^{\mathcal{L}}(x,y) = \sum \operatorname{Li}_{2}^{\mathcal{L}}$ 's, and  $\operatorname{Li}_{1,1,1}^{\mathcal{L}}(x,y,z) = \sum \operatorname{Li}_{3}^{\mathcal{L}}$ 's.

✓ (Intro)

Weight 4: After tweaking variables  $\overline{\Delta} \operatorname{Li}_{3,1}^{\mathcal{L}}(\frac{x}{y},y) = -\operatorname{Li}_{2}^{\mathcal{L}}(x) \wedge \operatorname{Li}_{3}^{\mathcal{L}}(y)$ .

Vanishing of  $\overline{\Delta}\left(\sum_{i}\lambda_{i}\operatorname{Li}_{3,1}^{\mathcal{L}}(\frac{x_{i}}{y},y)\right)$  under dilogarithm identities  $\sum_{i}\lambda_{i}\operatorname{Li}_{2}(x_{i})=0$ .

- Symmetries  $\operatorname{Li}_2^{\mathcal{L}}(x) + \operatorname{Li}_2^{\mathcal{L}}(1-x) = 0$  and  $\operatorname{Li}_2^{\mathcal{L}}(x) + \operatorname{Li}_2^{\mathcal{L}}(x^{-1}) = 0$  (Intro  $+\varepsilon$ )
- 5-term relation  $\sum_{i=0}^4 (-1)^i \operatorname{Li}_2^{\mathcal{L}}([z_0,\ldots,\widehat{z_i},\ldots,z_4]) = 0$ ,  $[z_1,\ldots,z_4] = \frac{(z_1-z_2)(z_3-z_4)}{(z_2-z_3)(z_4-z_1)}$

Theorem (5-term reduction, Gangl 2012; Goncharov-Rudenko, Matveiakin-Rudenko)

$$\sum_{i=0}^{4} (-1)^i \operatorname{Li}_{3,1}^{\mathcal{L}} \left( \frac{\operatorname{cr}(z_0, \dots, \widehat{z}_i, \dots, z_4)}{y}, y \right) = \sum_{i=1}^{122+\varepsilon} \operatorname{Li}_4^{\mathcal{L}} (g_i(z_0, \dots, z_4, y)), \quad g_i \text{ rational}$$

Remark: Crucial ingredient to Goncharov-Rudenko proof of Zagier's Conjecture on  $\zeta_F(4)$ .

# Iterating to understand higher depth

■ Computing  $\overline{\Delta}$  gets complicated:

$$\overline{\Delta} \operatorname{Li}_{4,1,1}^{\mathcal{L}}(x,y,z) = \operatorname{Li}_{2}^{\mathcal{L}}(y) \wedge \operatorname{Li}_{3,1}^{\mathcal{L}}(xy,z) - \operatorname{Li}_{3,1}^{\mathcal{L}}\left(yz,\frac{1}{y}\right) \wedge \operatorname{Li}_{2}^{\mathcal{L}}(xyz) - \operatorname{Li}_{3}^{\mathcal{L}}(y) \wedge \operatorname{Li}_{2,1}^{\mathcal{L}}(xy,z) \\ - \operatorname{Li}_{2,1}^{\mathcal{L}}\left(yz,\frac{1}{y}\right) \wedge \operatorname{Li}_{3}^{\mathcal{L}}(xyz) + \operatorname{Li}_{4}^{\mathcal{L}}\text{'s} \wedge \operatorname{weight} \ 2 + \operatorname{Li}_{n}^{\mathcal{L}}\text{'s} \wedge \operatorname{Li}_{m}^{\mathcal{L}}\text{'s}.$$

 $\blacksquare \ \, \mathsf{Use} \ \, \overline{\Delta} \ \, \mathsf{to} \ \, \mathsf{pick} \, \, \mathsf{"genuine"} \, \, \mathsf{depth} \, \, \mathsf{2} \, \, \mathsf{terms, so} \, \, \mathsf{iterate}^1 \, \, \, \overline{\Delta}^{[2]} \coloneqq \overline{\Delta} \circ (\overline{\Delta} \otimes \mathrm{id})$ 

$$\overline{\Delta}^{[2]} \operatorname{Li}_{4,1,1}^{\mathcal{L}}(x,y,z) = \operatorname{Li}_{2}^{\mathcal{L}}(y) \otimes \operatorname{Li}_{2}^{\mathcal{L}}(z) \otimes \operatorname{Li}_{2}^{\mathcal{L}}(xyz) \quad \text{(in CoLie_3)}.$$

Since  $\overline{\Delta}^{[2]} \operatorname{Li}_{a,b}^{\mathcal{L}}(x,y) = 0$ , we deduce  $\operatorname{Li}_{4,1,1}^{\mathcal{L}}(x,y,z) \neq \sum \operatorname{depth} 2$ .

Question: Is this the only obstruction?

#### Goncharov Depth Conjecture (simplified version)

A linear combination of MPL's has depth < k if and only if  $\overline{\Delta}^{[k-1]}$  vanishes.

<sup>&</sup>lt;sup>1</sup>needs to be formalised: in a coLie algebra, the coJacobi identity would imply naı̈e  $\Delta^2=0$ 

# Expectations and evidence (II)

Observe: If  $k > \frac{n}{2}$ , then  $\overline{\Delta}^{[k-1]}(\mathcal{L}_n) = 0$  (No space for k-many weight 2 pieces.)

#### Theorem (Rudenko, 2022)

Every weight n multiple polylogarithm can be expressed via depth  $\leq \lfloor \frac{n}{2} \rfloor$  functions.

Idea: Explicit formula via "quadrangular" polylogs with depth  $\frac{n}{2}$ 

Write:

$$f(x,y,z) := \operatorname{Li}_{4,1,1}^{\mathcal{L}}((xyz)^{-1}, y, z)$$

$$\overline{\Delta}^{[2]} f(x,y,z) = -\operatorname{Li}_{2}^{\mathcal{L}}(x) \otimes \operatorname{Li}_{2}^{\mathcal{L}}(y) \otimes \operatorname{Li}_{2}^{\mathcal{L}}(z) \quad (\text{in CoLie}_{k})$$

Then f(x,y,z) behaves like  $\mathrm{Li}_2^{\mathcal{L}}$  in each argument. Expect the following reductions:

$$f(x,y,z) + f(\underbrace{1-x}_{\text{or }x^{-1}},y,z) \stackrel{?}{\equiv} 0 \pmod{\mathsf{dp}} \leq 2)$$
 (Zagier-type)

### Weight 6 results

Saw similar expectations hold in weight 4. They are also recently known in weight 6.

Write 
$$f(x, y, z) := \operatorname{Li}_{4,1,1}^{\mathcal{L}}((xyz)^{-1}, y, z)$$
, as before

Theorem (Gangl-type reduction, Matveiakin-Rudenko 2022)

$$f(\text{dilog 5-term}, y, z) \equiv 0 \pmod{\text{depth}} \leq 2 \& \text{the Zagier reductions}$$

Theorem (Zagier-type reductions, C 2024, arXiv:2405.13853)

$$f(x, y, z) + f(1 - x, y, z) \equiv 0 \pmod{\text{depth}} \le 2$$
$$f(x, y, z) + f(x^{-1}, y, z) \equiv 0 \pmod{\text{depth}} \le 2$$

Consequence: Goncharov's Depth Conjecture holds for weight 6 depth 3 (k = 3).

(Zagier-reductions seem more difficult than the Gangl-reduction in general!)

Remark: Similar conjectural expectations in all even weights.

# Ingredients for the proof – Quadrangular polylogarithms

Rudenko's quadrangular polylogarithms  $\mathrm{QLi}_{n+k}(x_0,\ldots,x_{2n+1})$ 

- "Nice" MPL functions on  $\mathfrak{M}_{0,2n+2}$ , expressible via "quadrangulations" of polygons
- Weight-independent depth n expression via cross-ratio  $[x_a, x_b, x_c, x_d] = \frac{(x_a x_b)(x_c x_d)}{(x_b x_c)(x_d x_a)}$

$$QLi_{k+2}(x_{1},...,x_{6}) = Li_{k;1,1}^{\mathcal{L}} \left( + \underbrace{x_{1}^{2} x_{1}^{2} x_{1}^{2}}_{x_{5}} - \underbrace{x_{2}^{2} x_{1}^{2} x_{1}^{2}}_{x_{5}} + \underbrace{x_{2}^{2} x_{1}^{2}}_{x_{5}} + \underbrace{x_{3}^{2} x_{1}^{2}}_{x_{5}} \right)$$

$$= Li_{k;1,1}^{\mathcal{L}} \left( [x_{1}, x_{2}, x_{3}, x_{6}], [x_{3}, x_{4}, x_{5}, x_{6}] \right) - Li_{k;1,1}^{\mathcal{L}} \left( [x_{1}, x_{2}, x_{5}, x_{6}], [x_{3}, x_{4}, x_{5}, x_{2}] \right)$$

$$+ Li_{k;1,1}^{\mathcal{L}} \left( [x_{1}, x_{4}, x_{5}, x_{6}], [x_{1}, x_{2}, x_{3}, x_{4}] \right)$$

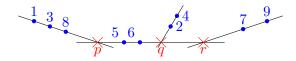
#### General functional equation

$$\mathbf{Q}_{2k}$$
:  $\sum_{i=1}^{2k+3} (-1)^i \operatorname{QLi}_{2k}(x_1, \dots, \widehat{x_i}, \dots, x_{2k+3}) \equiv 0 \pmod{\operatorname{depth}} < k$ 

# Ingredients (II) – Stable genus 0 curves

Part of boundary strata of Deligne-Mumford compactification  $\overline{\mathfrak{M}}_{0,n}$ 

- lacksquare Components are isomorphic to  $\mathbb{P}^1$
- Only singular points are simple double points
- Number of marked & singular points per components  $\geq 3$



Gadget to encode limits of projective configurations

Idea: The points  $x_1, x_3, x_8$  all go to p

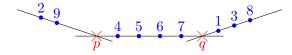
But: There is a always projective transformation moving  $(x_1, x_3, x_8, p)$  to  $(\infty, 0, 1, z)$ 

So: Points  $x_1, x_3, x_8$  split off as a separate  $\mathbb{P}^1$ 

Cross-ratios well-defined:  $[x_1, x_3, x_8, x_i] = [x_1, x_3, x_8, p]$  and  $[x_1, x_3, x_2, x_4] = [x_1, x_3, p, p] = 0$ 

# Ingredients (III) – Build to symmetries step-by-step

Specialise  $\mathbf{Q}_6$  to



and use  $\lim_{x_i \to 0} \operatorname{or}_{\infty} \operatorname{Li}_{k_1, \dots, k_d}^{\mathcal{L}}(x_1, \dots, x_d) = 0 \pmod{\mathsf{dp}} < d$ 

Write  $f(x, y, z) := \operatorname{Li}_{4,1,1}^{\mathcal{L}}((xyz)^{-1}, y, z)$ , as before

#### Lemma (Full Symmetry 1)

$$f(A, B, C) \equiv -f\left(1 - A, \frac{B}{B - 1}, 1 - C\right) \pmod{\text{depth 2}}$$

Remark: First needed to show  $f(1, x, y) \equiv 0 \pmod{\text{depth 2}}$ , already non-trivial(!)

- Find a second *different* symmetry, and a four-term relation
- Play the symmetries against each other
- Extract (eventually...) the Zagier-type reduction

- Depth reductions of MPL's
  - Appearances in Number Theory, Hyperbolic Geometry, High-Energy Physics
- Goncharov's Depth Conjecture
  - Cobracket on (motivic) MPL's
  - Truncation and iteration of the cobracket
  - Cobracket conjecturally detects depth
- Zagier-, & Gangl-type reductions in weight 4, and weight 6
- Ideas from the proof
  - Quadrangular polylogarithm functional equations
  - Degenerations to stable genus 0 curves
  - Discover and exploit many symmetries/short identities